# Tree-like width of knots and obstructions. 

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## Knots

Knot:
A knot is a tame embedding $\mathbb{S}^{1} \rightarrow \mathbb{S}^{3}$ considered up to continuous deformation (ambient isotopy).


Natural algorithmic questions:

- Is a given knot equivalent to the trivial one ? (NP and co-NP)


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- Are two given knots equivalent ?


## "Width" invariant of knots

Aim to measure some notion of "width" on knots.

## Examples of width invariants :

- Trunk: Minimal number among linear sweepouts of the maximal number of intersections between the knot and a plane (or a sphere).

- Bridge number: Minimal number of trivial arcs on a sphere that "splits" the knot.


## Treewidth on knots

Aim to measure "how close" a knot is to a tree.


A diagram of a knot is a 4-valent graph.
Diagrammatic treewidth of a knot $K$ :
It is the minimum treewidth of a diagram of $K$.

## Treewidth of a graph

Aim to measure "how close" a graph is to a tree.
Treewidth of a graph $G$ :
Maximum size of a bag in a tree decomposition of $G$.

In a tree decomposition :

- Each vertex is in a bag
- Each pair of edge extremities is in a bag
- Bags containing a fixed vertex create a subtree



## Treewidth of a graph

Small treewidth :


High treewidth :


## Why is treewidth interesting?

- Key tool for dynamic programming on graphs. For example: Jones polynomial, Kauffman polynomial [J. A. Makowsky and J. P. Mariño, 2003], HOMFLY-PT polynomial [B. A. Burton, 2018], and quantum invariants [C. Maria, 2019] are efficiently computable on small treewidth diagrams of knots.


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- It plays a major role in the graph minor theory from Robertson and Seymour.
- Occurs naturally in other areas:

For example: 3-manifolds [K. Huszár and J. Spreer, 2019], lower bounds distortion on knots...

## Back to knots

Knots always have diagrams with high treewidth.


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Knots always have diagrams with high treewidth.


Question from [J. A. Makowsky and J. P. Mariño., The parametrized complexity of knot polynomials, 2003] :
Are there knots for which all diagrams have high treewidth ?

## Back to knots

Theorem [A. De Mesmay, J. Purcell, S. Schleimer, and E. Sedgwick, On the tree-width of knot diagrams., 2019]:
Let $T_{p, q}$ be a torus knot. Then $\operatorname{tw}\left(T_{p, q}\right)=\Omega(\min (p, q))$.


Towards a 3D definition:
The proof uses a width notion generalising carving width (treewidth equivalent) on knots, using multiple Heegaard splittings.

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Towards a 3D definition:
The proof uses a width notion generalising carving width (treewidth equivalent) on knots, using multiple Heegaard splittings.
My goal:
Define a width invariant on knots generalising branchwidth: definition stemming from structural graph theory tailored to proving lower bounds.

## Branchwidth definition

Branchwidth is another equivalent of treewidth : $\frac{2}{3} t w \leq b w \leq t w$. Branchwidth of a Graph :
$T$ is a trivalent tree and $\phi: E(G) \rightarrow L(T)$ is a bijection.

$$
\operatorname{bw}(G)=\min _{(T, \phi) \in \operatorname{BD}(G)} \max _{e \in T}\left|\partial \phi^{-1}\left(L\left(T_{1}\right)\right)\right|
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Branchwidth can be lower bounded through its dual problem: existence of a tangle.

## Tangle on graphs

## Tangle definition

A tangle of order $k$, noted $\mathcal{T} \subset \mathcal{P}(E)$ is a collection of subsets of $E$ ("small sides") such that:

- $\forall A \in \mathcal{T}, \partial A<k$, sets in the tangle have boundary less than $k$.
- For all $A, B \in \mathcal{P}(E)^{2}$, if $(A, B)$ is a separation of order less than $k, \mathcal{T}$ contains $A$ or $B$.
- $\forall A \in \mathcal{P}(E), \forall B \in \mathcal{T}, A \subset B$ and $\partial A<k \Rightarrow A \in \mathcal{T}$. Small sides are consistent.
- $\forall A, B, C \in \mathcal{P}(E)^{3}, A \sqcup B \sqcup C=E$ $\Rightarrow\{A, B, C\} \not \subset \mathcal{T}$. No three small sides can cover the whole space.

- $\forall e \in E, E \backslash\{e\} \notin \mathcal{T}$. A single edge is small.


## Obstruction



A branch decomposition of width 3 implies the impossibility for a tangle of order 3.

## Branchwidth definition



A double bubble : two spheres that intersect on a single disk.

## Branchwidth definition



A branch-decomposition or sweep of $\mathbb{S}^{3}$ is a continuous map $f: \mathbb{S}^{3} \rightarrow T$ where $T$ is a trivalent tree such that :

$$
f^{-1}:\left\{\begin{array}{cl}
\text { leaf } & \mapsto \text { point } \\
\text { vertex } & \mapsto \text { double bubble } \\
\text { point interior to an edge } & \mapsto \text { sphere }
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$$

Branchwidth
The branchwidth of $K$ written $\mathbf{b w}(K)$ is:

$$
\operatorname{bw}(K)=\inf _{f} \sup _{e \in E(T), x \in \stackrel{e}{e}}\left|f^{-1}(x) \cap K\right| .
$$

## Bubble tangle on knots

Bubble tangle
Let $K$ be a knot, tangle of order $n$, noted $\mathcal{T}$ is a collection of closed balls of $\mathbb{S}^{3}$ such that:

- All subsets of $\mathcal{T}$ have less than $n$ intersections with $K$ on their boundaries.
- For any spheres $S$ of $\mathbb{S}^{3}$, if $|S \cap K|<n$ then a component of $\mathbb{S}^{3} \backslash S$ is in $\mathcal{T}$.
- For any three closed balls $B_{1}, B_{2}, B_{3}$ that induces a double bubble, not all three of $B_{1}, B_{2}, B_{3}$ are in $\mathcal{T}$.
- $T$ contains trivial balls:



## Results

Theorem 1
If there exists a bubble tangle of order $n$ on $K$ then $b w(K) \geq n$.

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Theorem 2
Let $n$ be the maximal order of a bubble tangle on $K$. Then $\operatorname{bw}(K)=n$.

## Distortion and representativity

## Representativity

If $K \hookrightarrow \Sigma \hookrightarrow \mathbb{S}^{3}$, the representativity of $K$ on $\Sigma$ is the minimum number of intersection between a non contractible, compressible curve on $\Sigma$ and $K$.


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## Representativity

If $K \hookrightarrow \Sigma \hookrightarrow \mathbb{S}^{3}$, the representativity of $K$ on $\Sigma$ is the minimum number of intersection between a non contractible, compressible curve on $\Sigma$ and $K$.


## Distortion

The distortion of $K \hookrightarrow \mathbb{R}^{3}$ is defined as $\delta(K)=\sup _{x, y \in K} \frac{d_{K}(x, y)}{d_{\mathbb{R}^{3}}(x, y)}$.
Pardon showed that representativity of torus knots lower bound their distortion [Pardon, 2011].

## Bubble tangle from representativity

Generalising the work of [Pardon, 2011],
Theorem 3
Let $K$ be a knot of representativity $n$ on a surface $\Sigma$. Then there exist a bubble tangle of order $c \times n$ (where $\left.c \in\left[\frac{1}{2}, \frac{4}{3}\right]\right)$.


Corollary
Torus knots have high treewidth.

## Conclusion

## Results:

- Branchwidth is a width invariant that can be defined on knots and generalises the notion from structural graph theory.
- It can be easily lower bounded through the dual problem: bubble tangle.
- These obstructions can arise from representativity.
- Extends naturally to spatial graphs


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Thank you for listening !


Removing an inessential curve from $S$.


## Transformation from a double bubble to $S^{\prime}$



Evolution of two circles of $S_{t} \cap M$ when $t$ grows.

## Diagram


$T_{9,7}$ and one of its graph.

